

" A txn T reads from a txn T' in a schedule S if some item X is first written by T' and later read by T . In addition T' should not have been aborted before T reads item X , and there should be no txns that write X after T' writes it and T reads it. (unless those txns, if any, have aborted before T reads X) "

T_1	T_2
$\{$ $w(X)$	
$\{$ aborted	$\{$
	$\{$ $r(X)$

Here, if T_1 aborts, then T_1 must be rolled back. So theoretically, T_2 reads the rolled back value of X , i.e, the value of X before T_1 made the write.

Hence, we cannot say in this case that T_2 reads from T_1 .
In other words, if T_2 reads from T_1 , such a scenario must not occur.